

Compressing Vector OLE

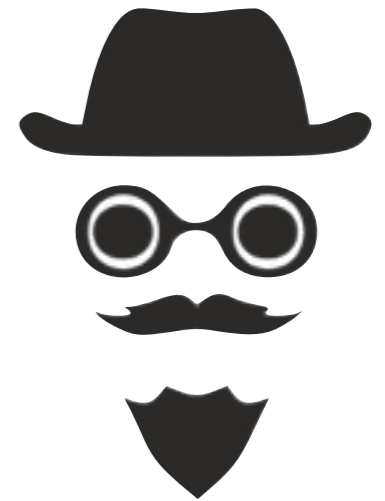
Elette Boyle, *Geoffroy Couteau*, Niv Gilboa, Yuval Ishai



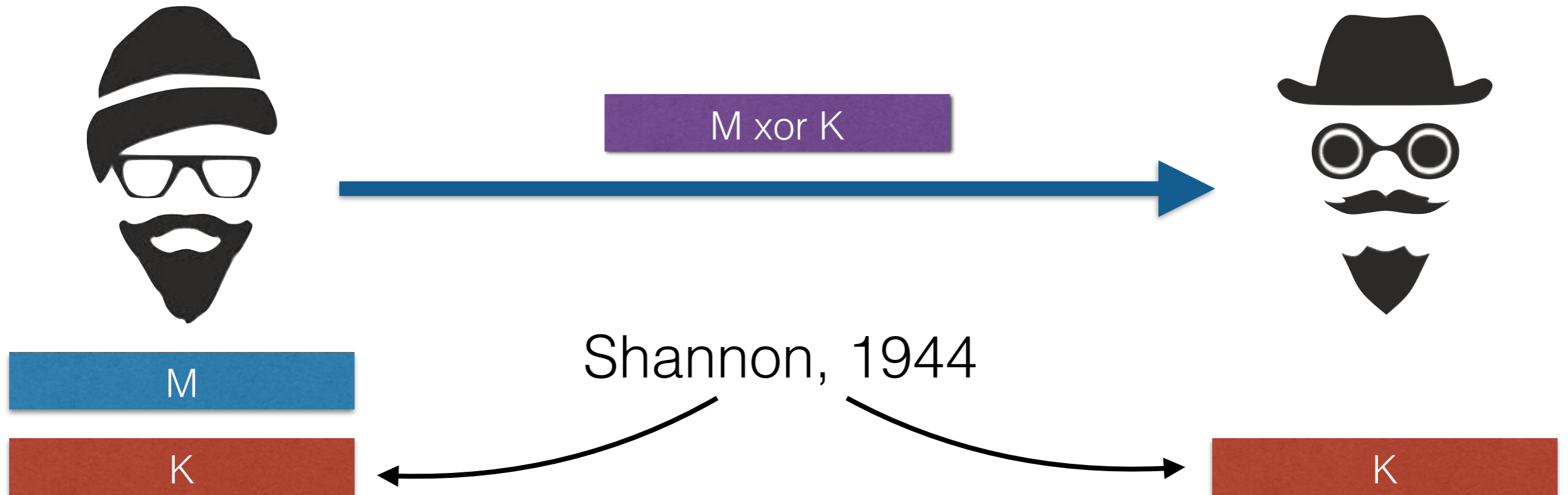
Secure Communication



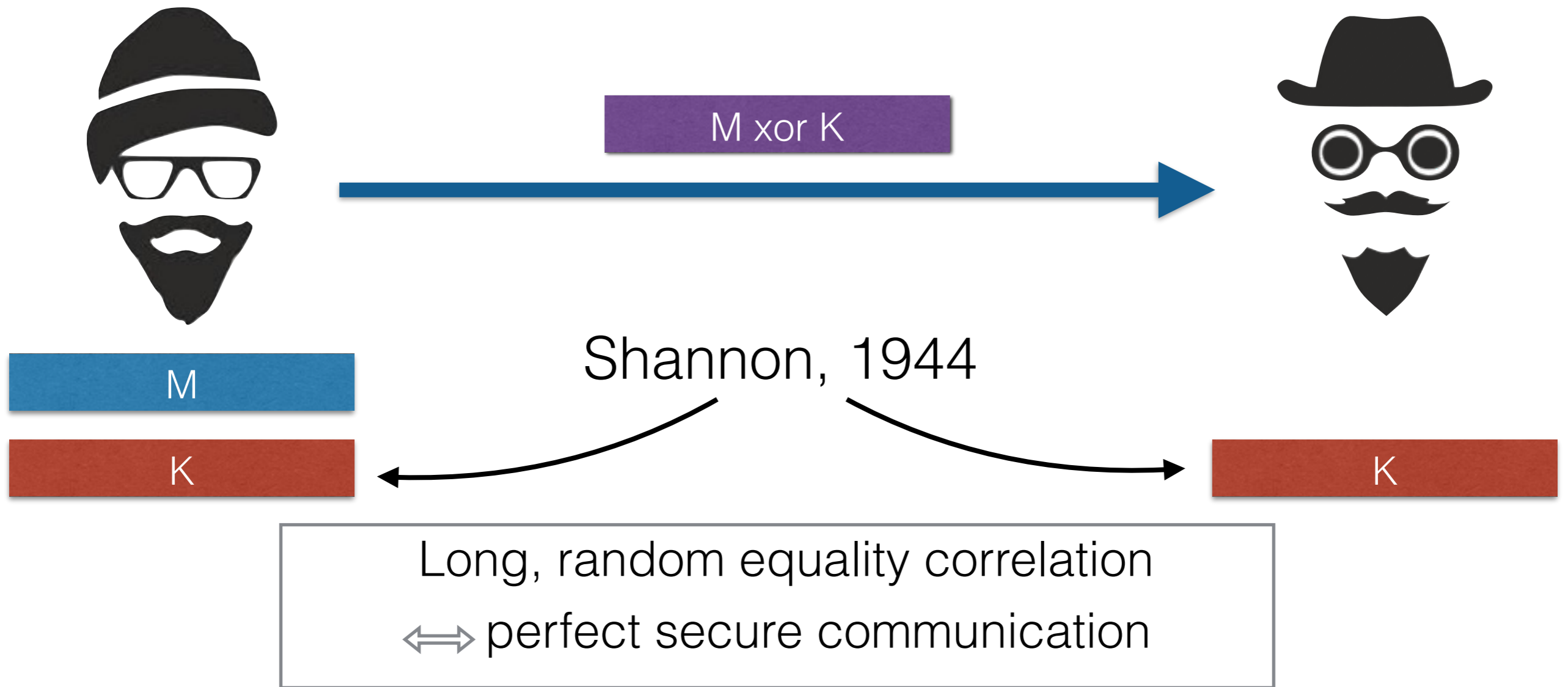
M



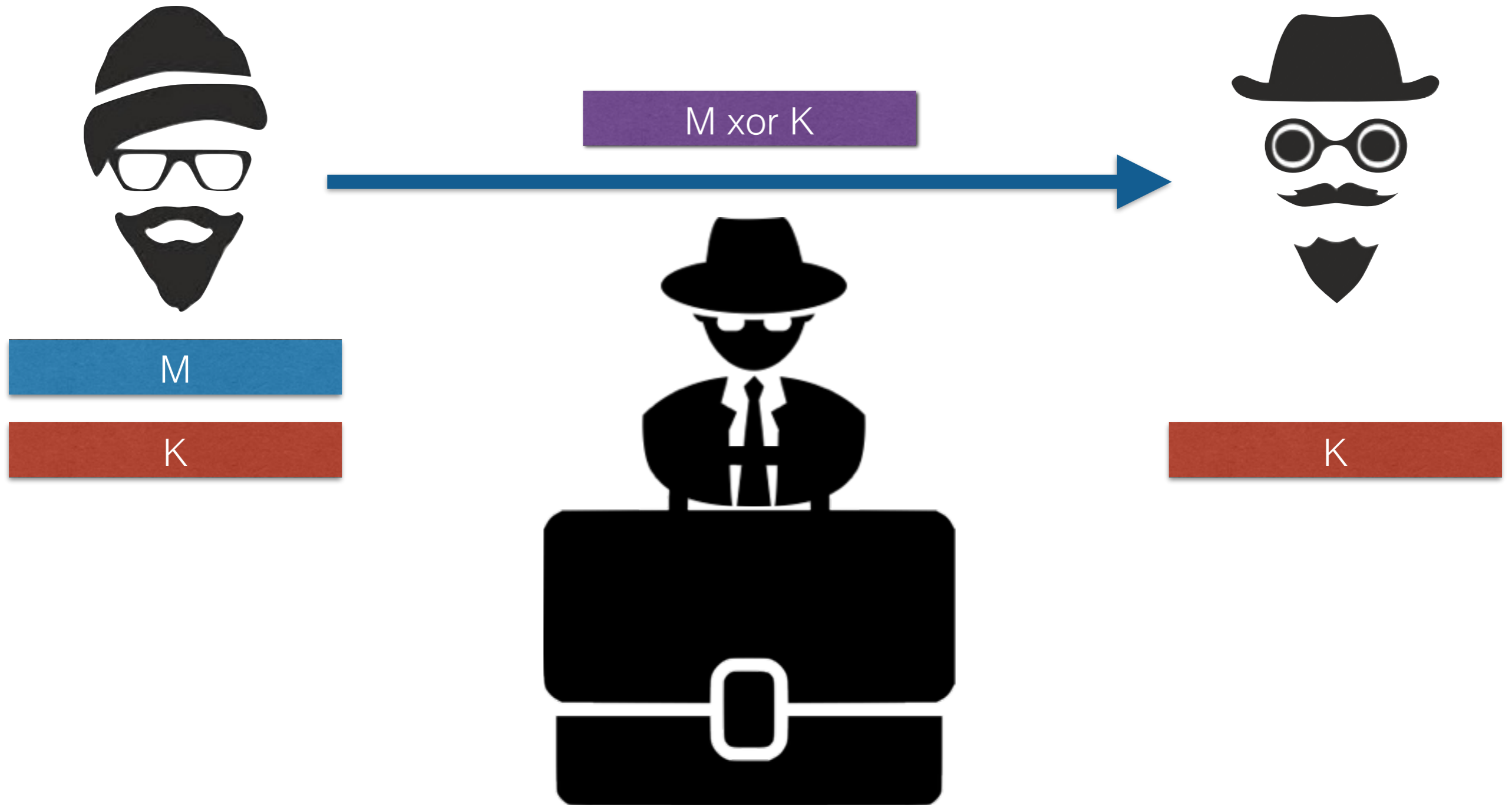
Secure Communication



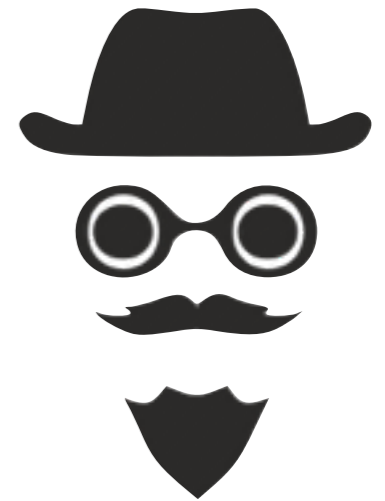
Secure Communication



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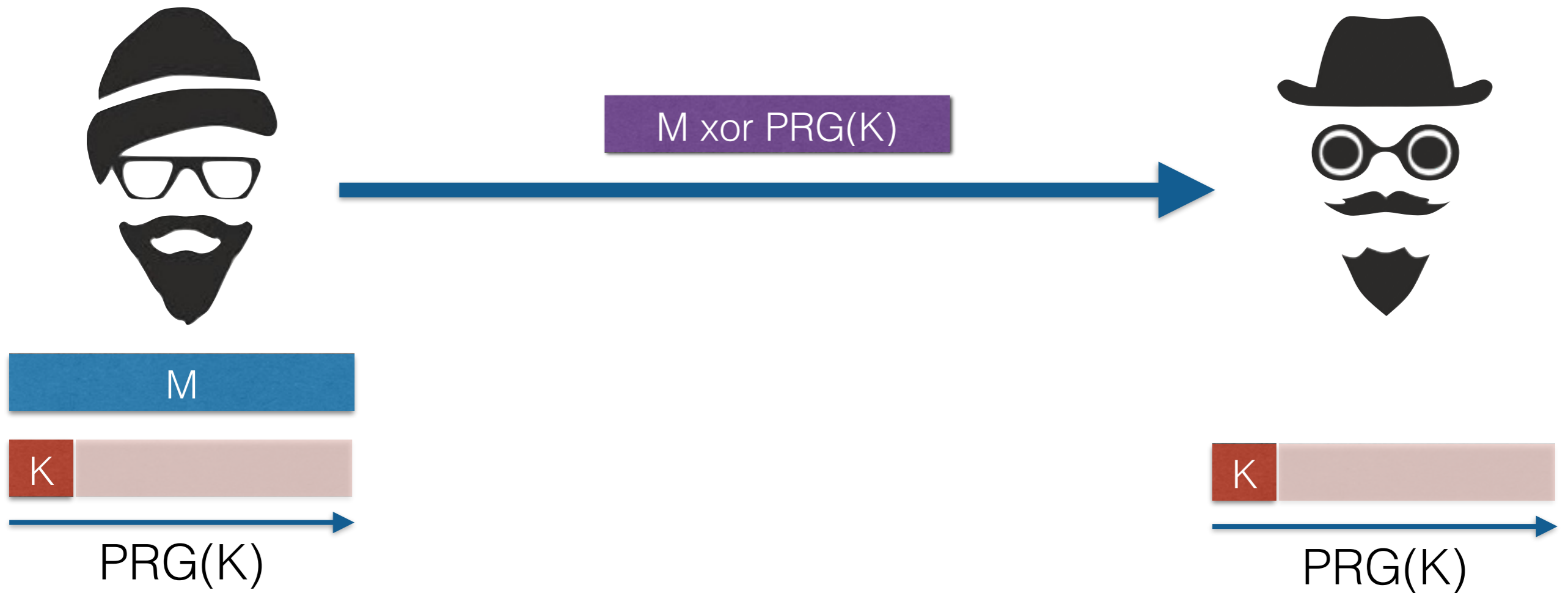


M

K

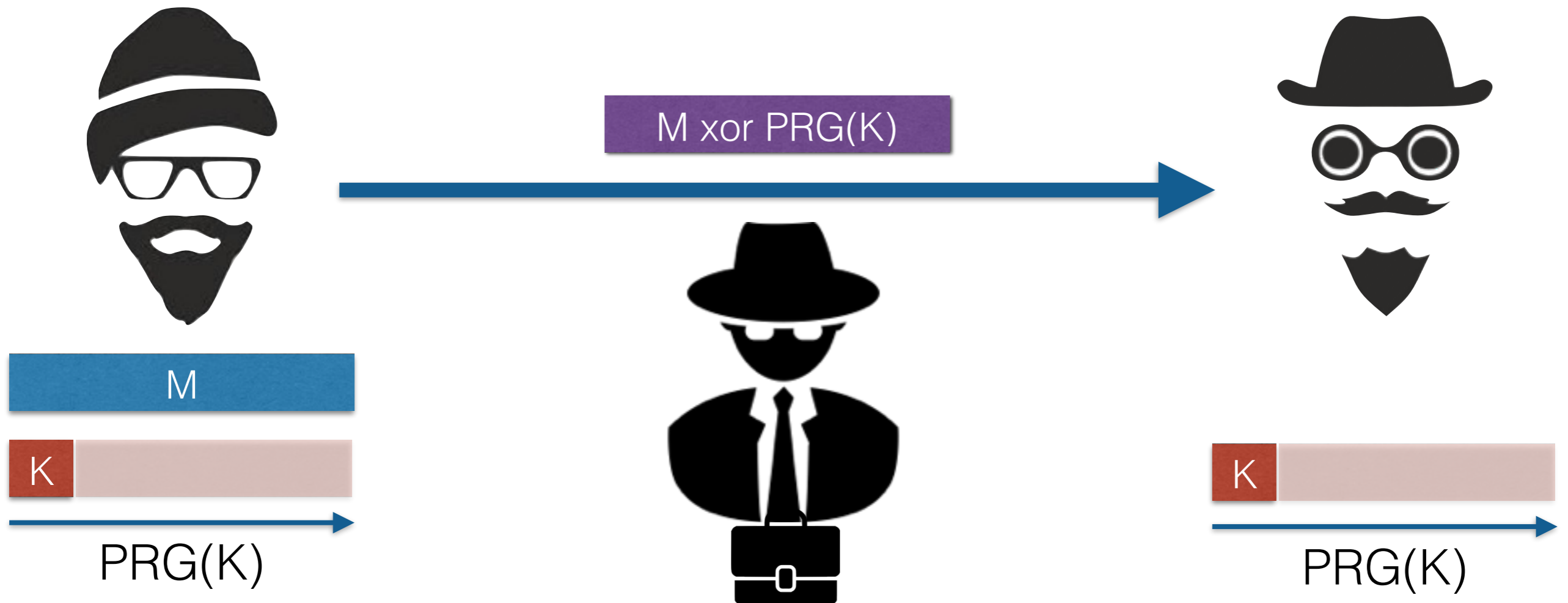
K

Secure Communication



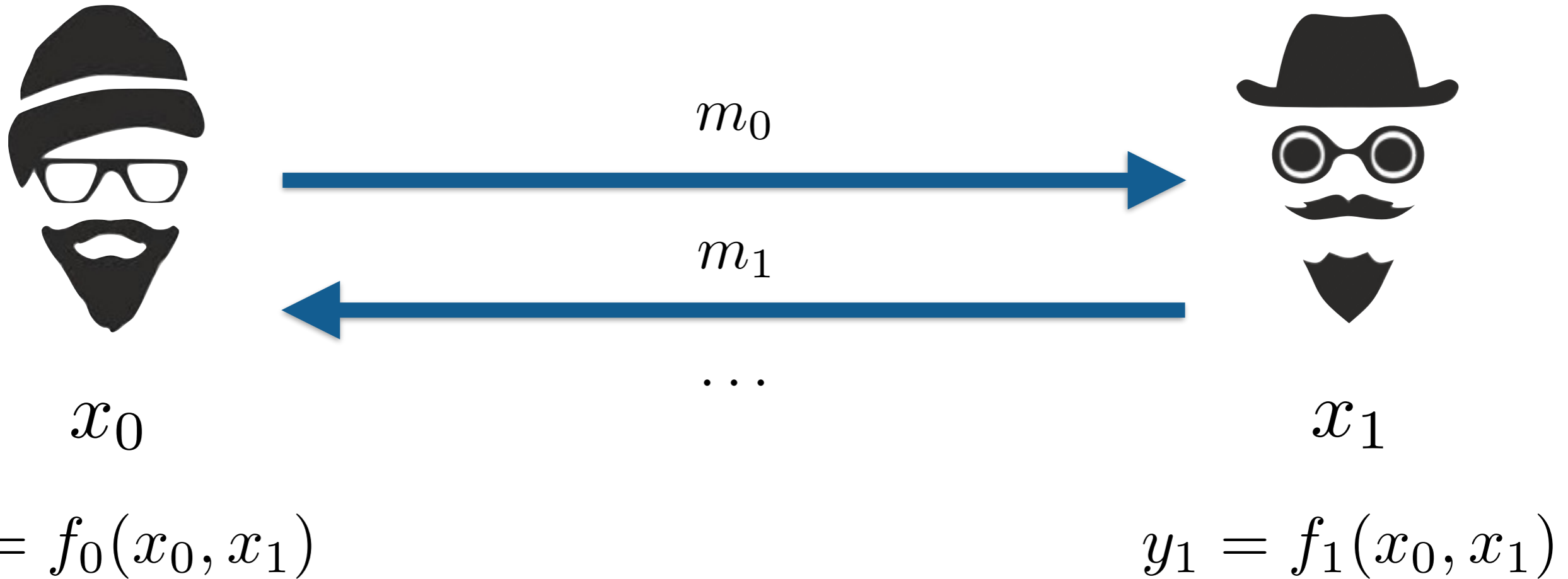
Short keys generate long (pseudo)-random correlations, enabling secure communication

Secure Communication



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Secure Computation



Secure Computation



x_0

$$y_0 = f_0(x_0, x_1)$$

K0

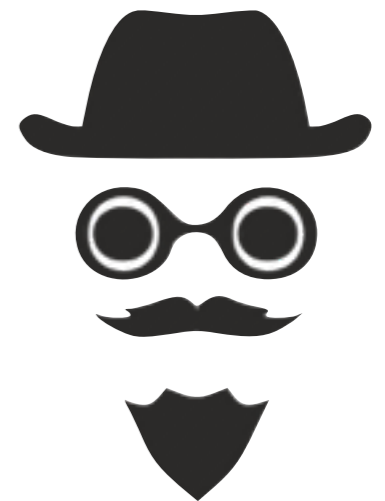
m_0



m_1



...



x_1

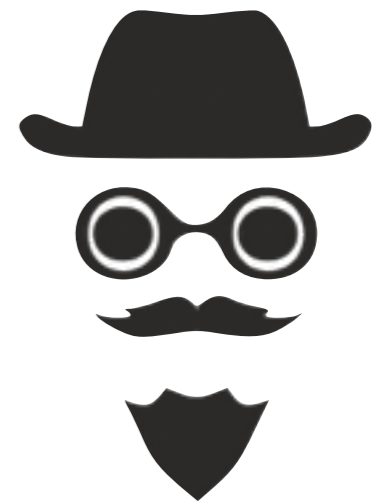
$$y_1 = f_1(x_0, x_1)$$

K1

Secure Computation



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x_1

m_0

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...

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Beaver, 1995

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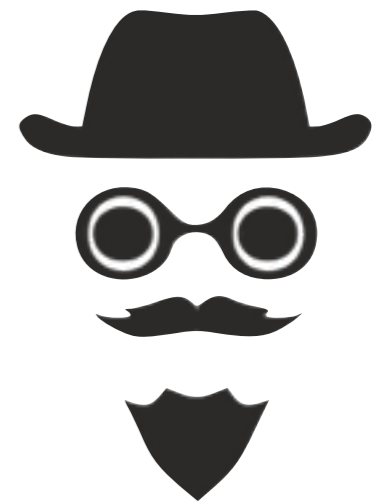
K1

Long, random 'more complex' correlation
↔ perfect secure computation

Secure Computation



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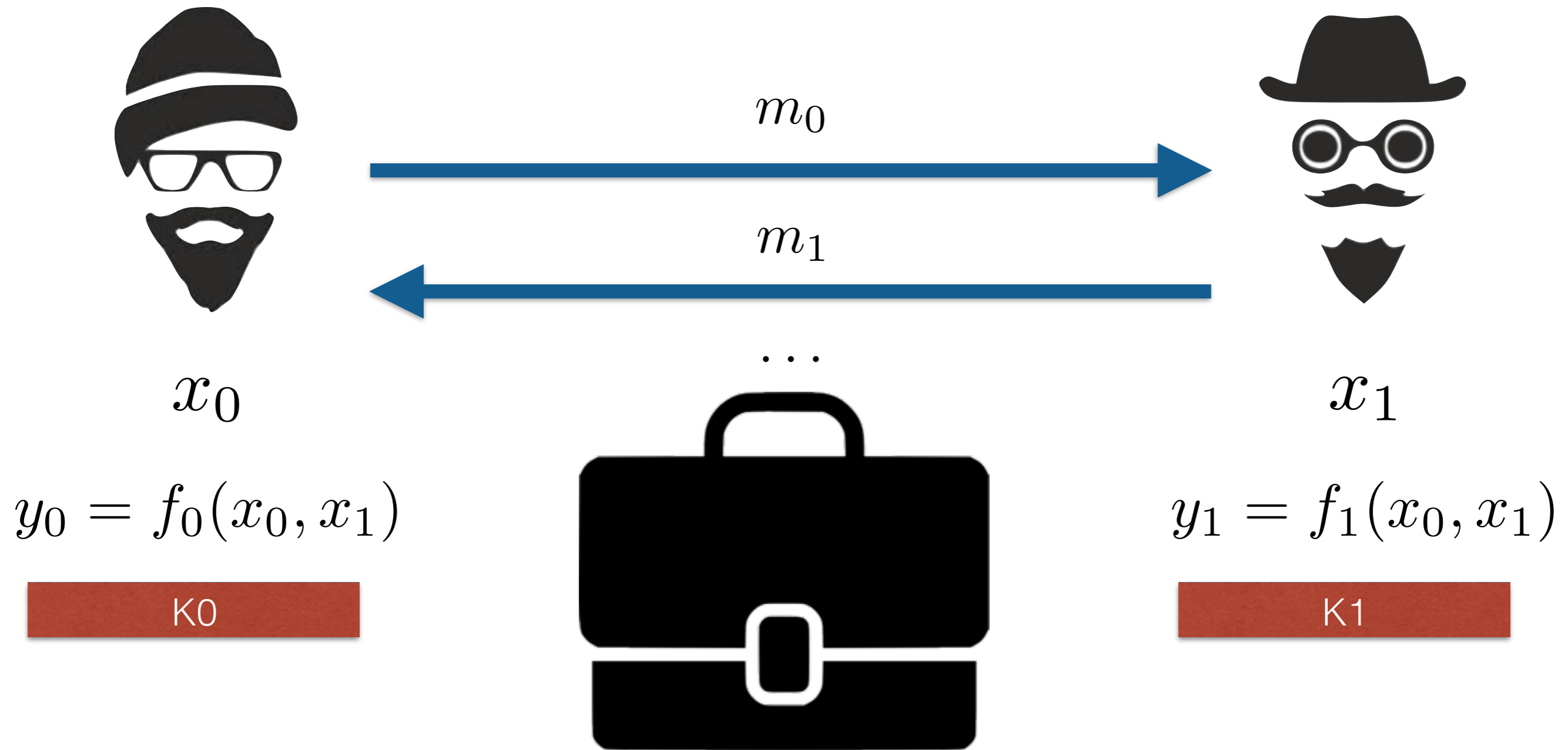
K0

K1

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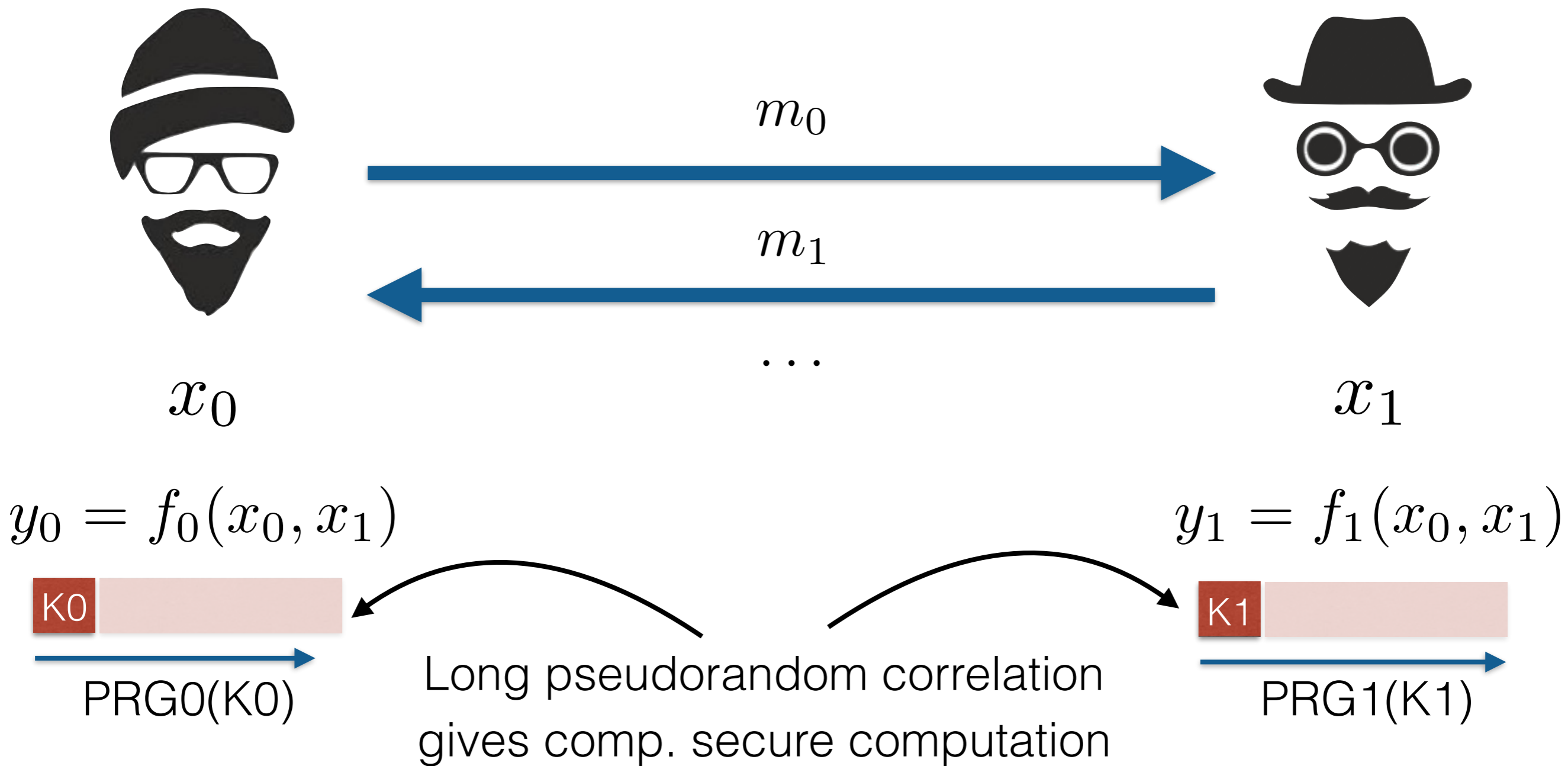
[Bea95, Bea97, IPS09, BDOZ11, BIKO12, NNOB12, DPSZ12, IKMOP13, DZ13, DLT14, BIKK14, LOS14, FKOS15, DZ16, KOS16, DNNR17, C18...]

Secure Computation

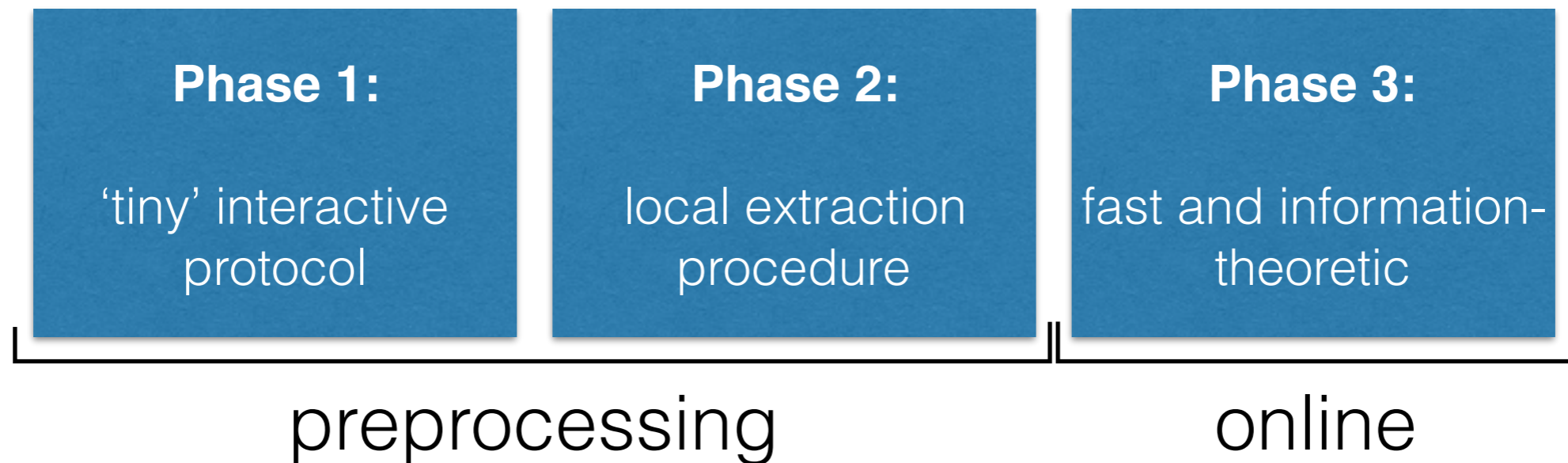


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Higher Degree Pseudorandom Correlations



Why would it be so cool?



- Reduced overall communication
- 'On demand' (or 'business-card') MPC
- Hiding communication pattern

Previous Works

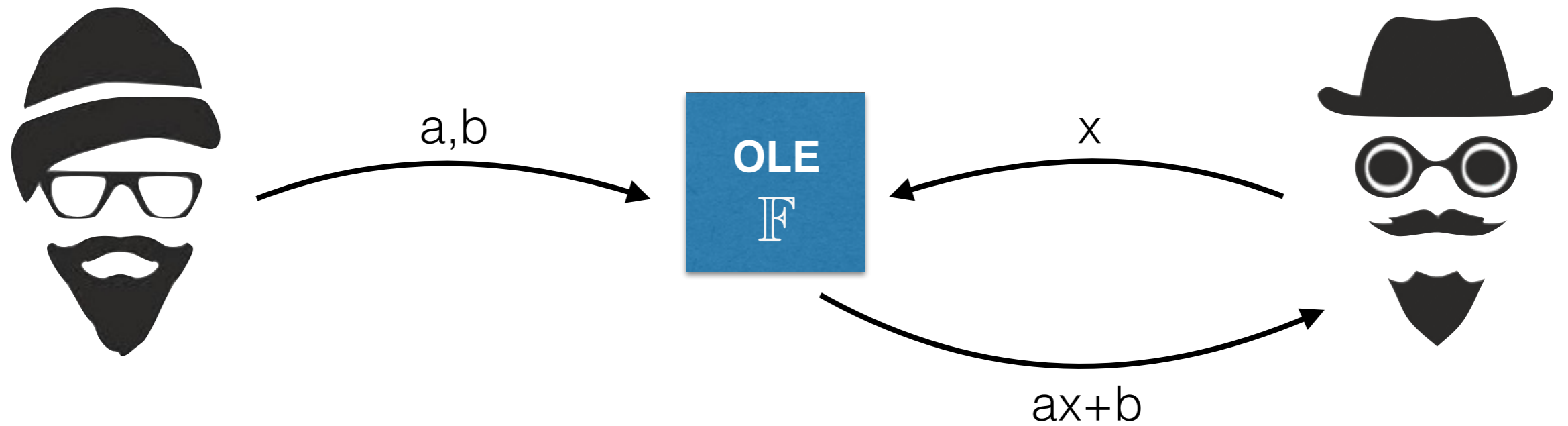
- GI99: multiparty linear correlations
- CDI05: generalization (pseudorandom Shamir)
- HIJKR16: arbitrary correlations, from iO
- BCGIO17: additive correlations, from HSS
- Scho18: relaxed form of additive correlations, key-homomorphic PRF

Previous Works

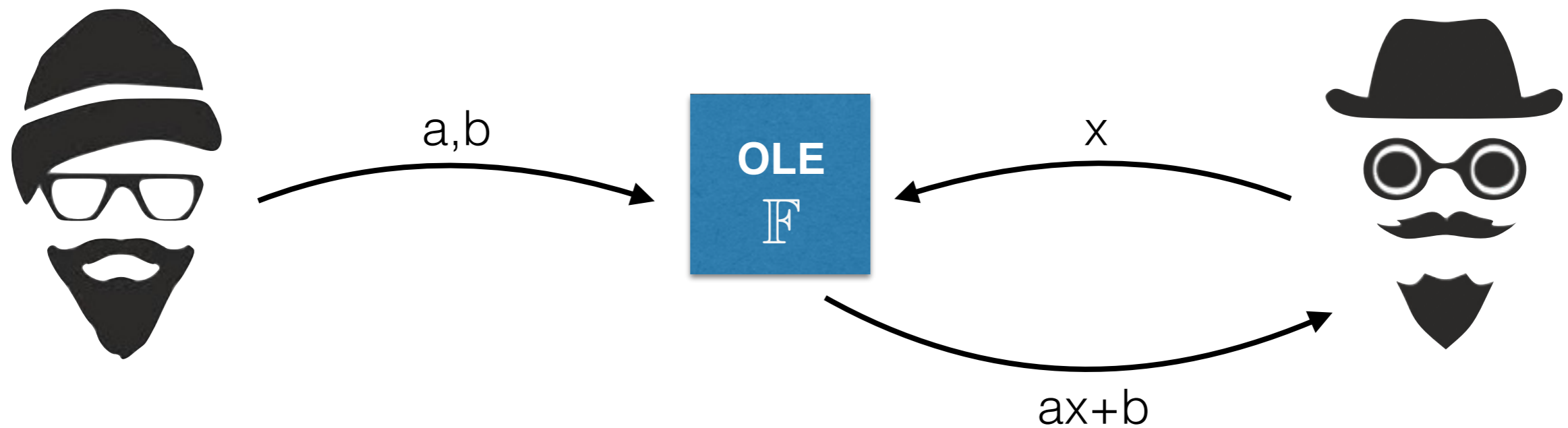
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Non-linear, 2-party: no efficient solution

Oblivious Linear Evaluation

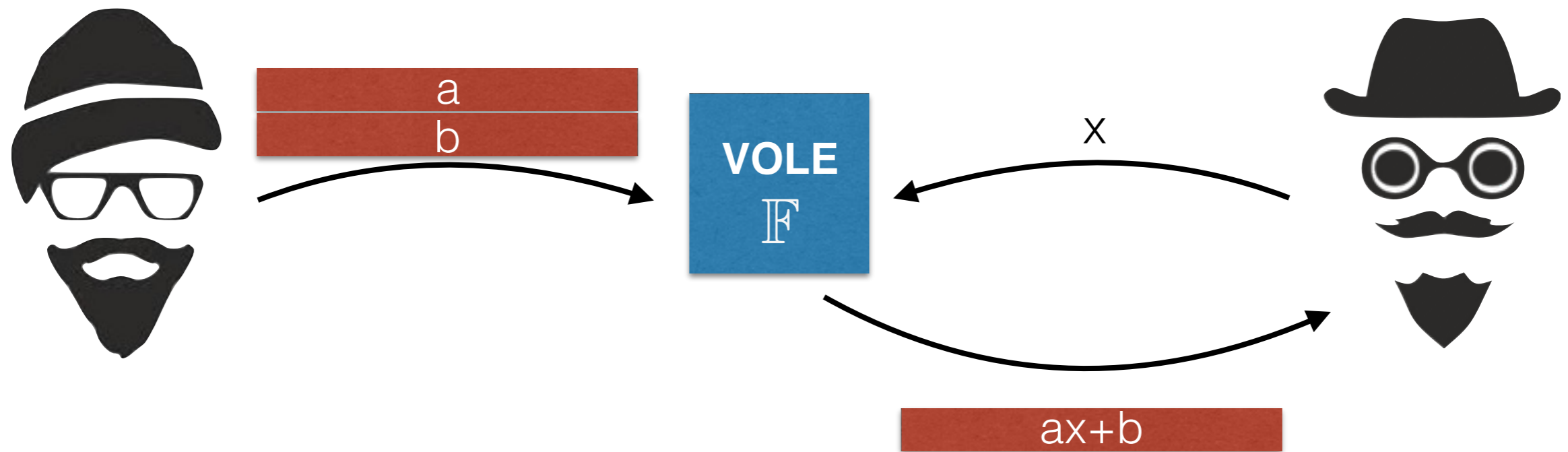


Oblivious Linear Evaluation

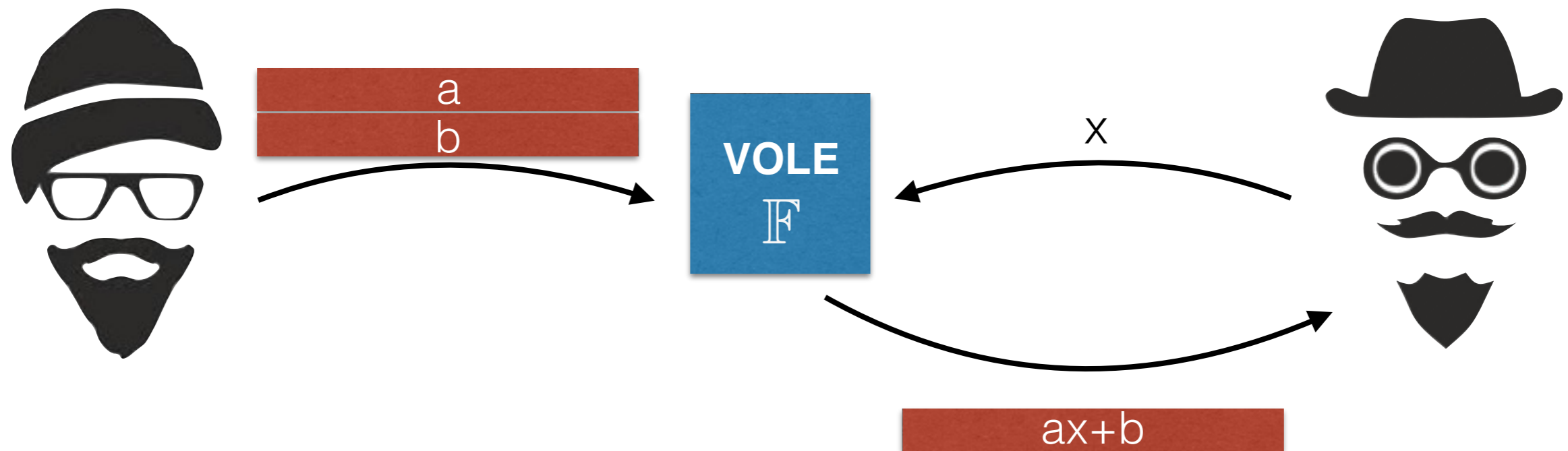


- Natural generalization of OT to large fields
- Enables secure computation of arithmetic circuits
- 'Dream goal' of correlation compression

Vector Oblivious Linear Evaluation



Vector Oblivious Linear Evaluation

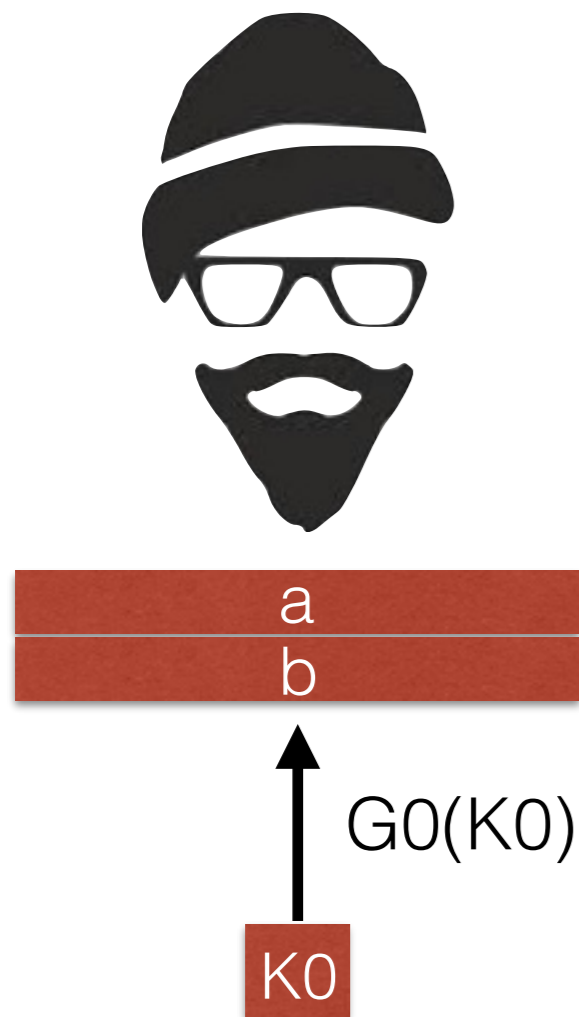


- Generalizes OLE the same way string-OT generalizes OT
- Still enjoys many cool applications for secure computation (more details later)

This Work: Efficient Compression for Vector-OLE Correlation



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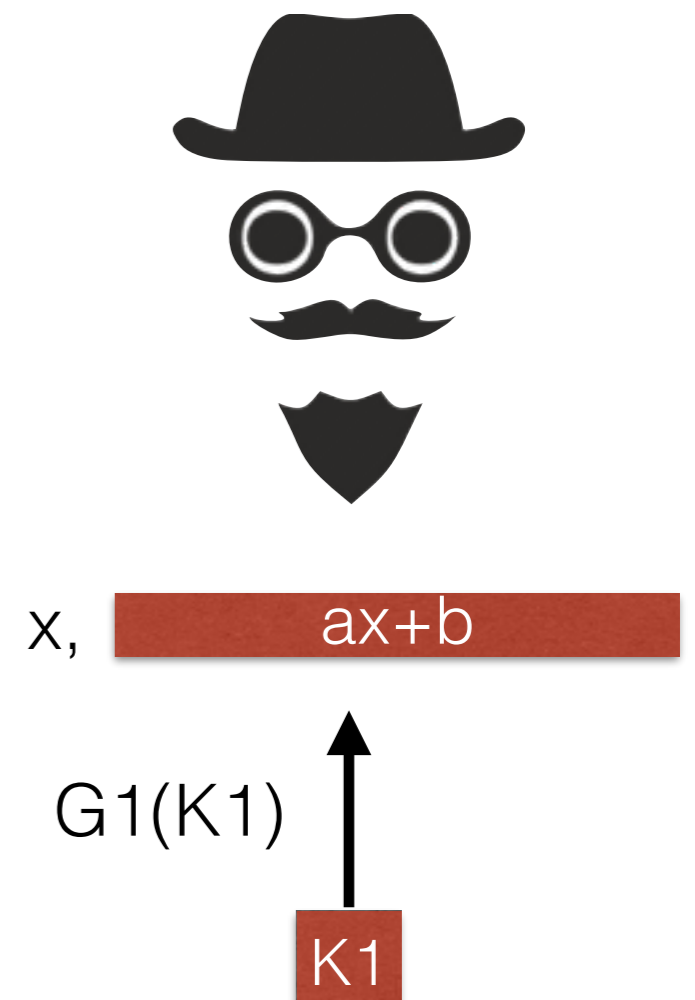


Security:

- Given K_1 , cannot distinguish (a,b) from random subject to $G_1(K_1) = ax+b$
- Given K_0 , learn nothing about x

Our Result:

An efficient* VOLE generator over any field under LPN



Applications

‘Direct’ applications:

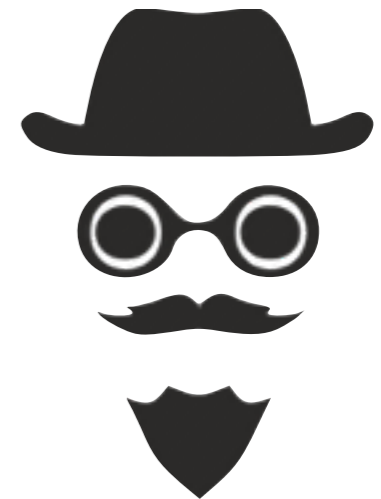
- Rate-1/2 Vector-OLE (better rate from LPN would be a breakthrough)
- Large fan-out MPC
 - Matrix-vector multiplication
 - Nearest neighbor search

‘High-end’ applications:

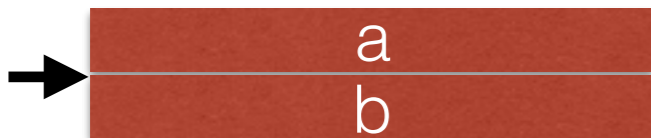
- Constant-depth arithmetic computation via arithmetic garbling
- Reusable non-interactive zero-knowledge in the preprocessing model

Our Method

1. *Sparse* VOLE-generator $\xrightarrow{\text{LPN}}$ pseudorandom VOLE-generator



K0



additive shares of $a \cdot x$

$x,$



K1

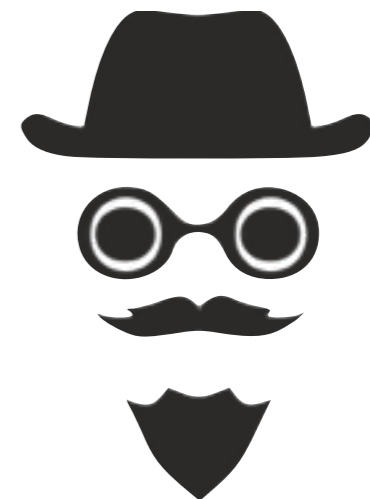
Our Method

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a

$0000 \dots a_i \dots 00$



K_0

a
 b

$x,$

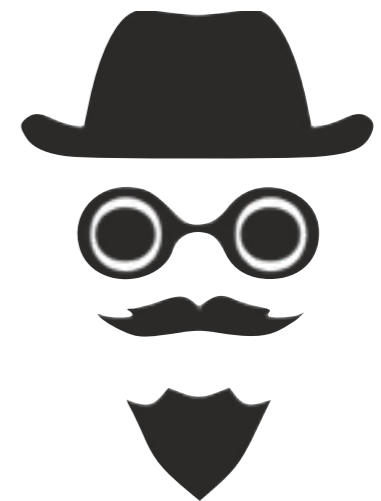
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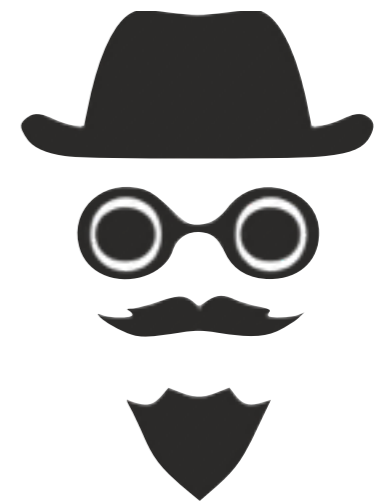
$$= \begin{matrix} \text{a} \\ 0000 \dots a_i \dots 00 \end{matrix}$$



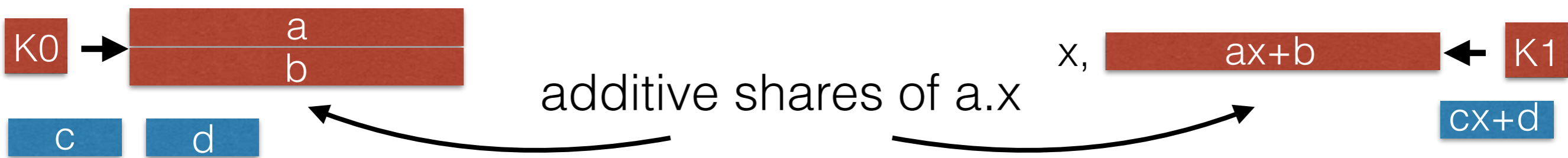
LPN: secret M + sparse error \approx random

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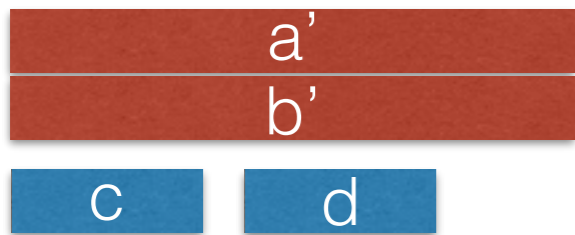


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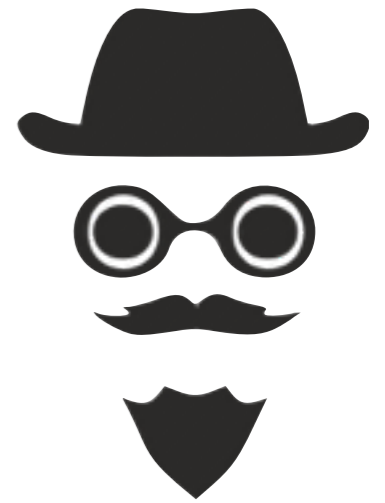


LPN: **secret** M + sparse error \approx random

Our Method



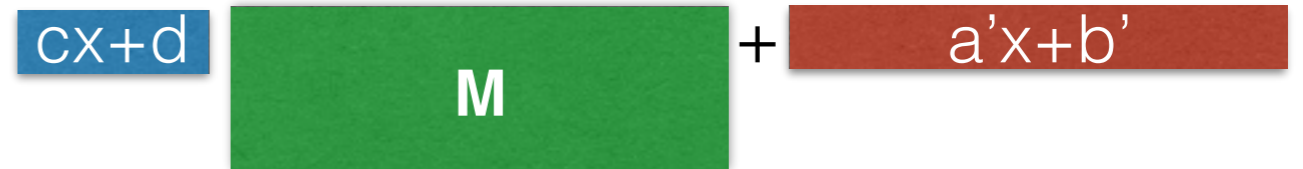
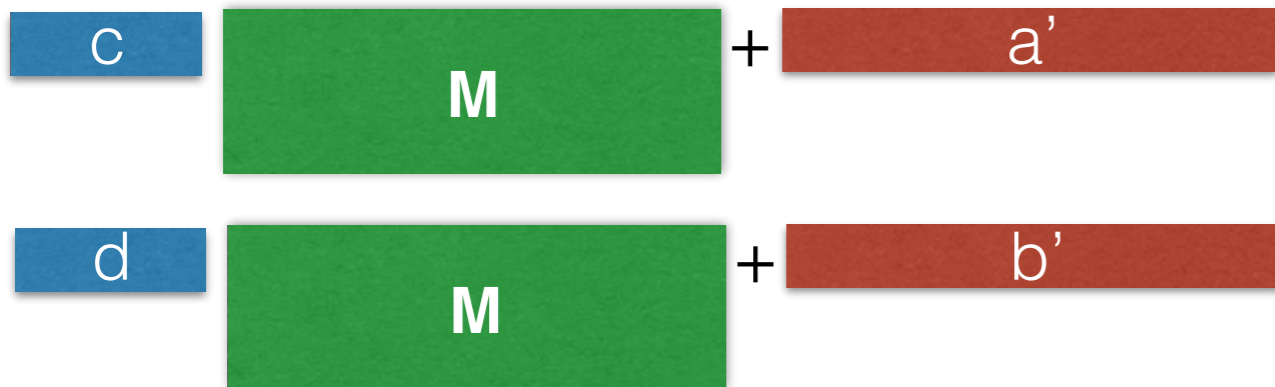
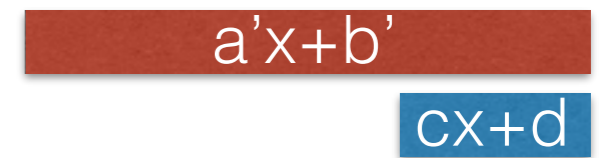
K0



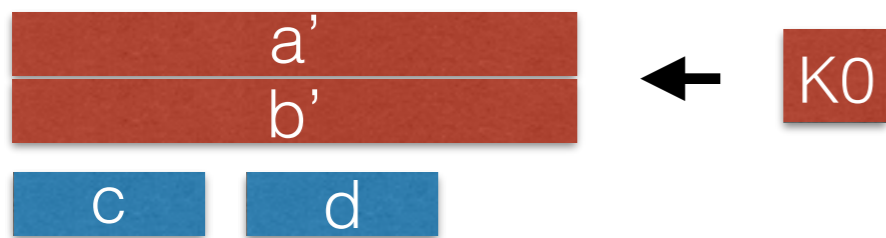
K1



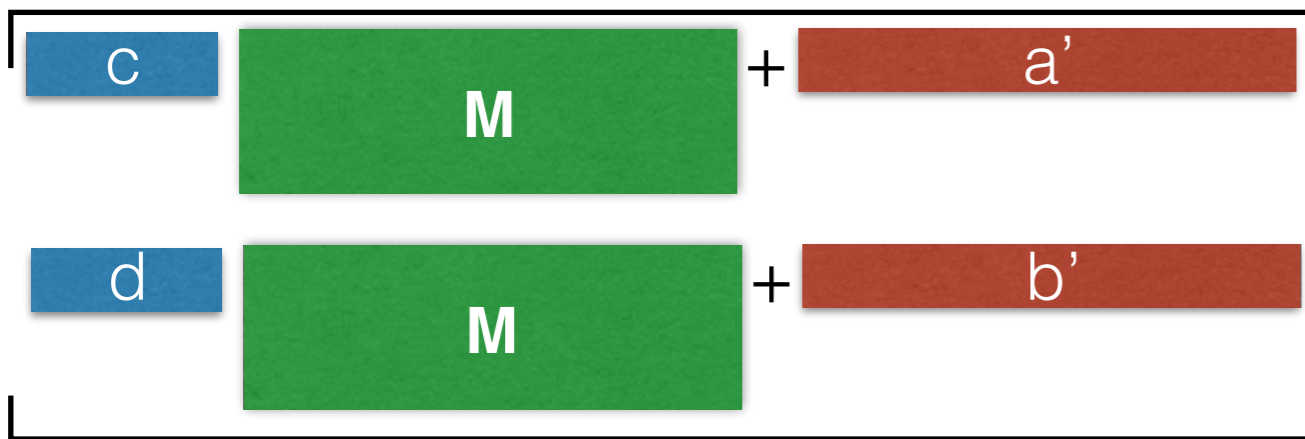
x,



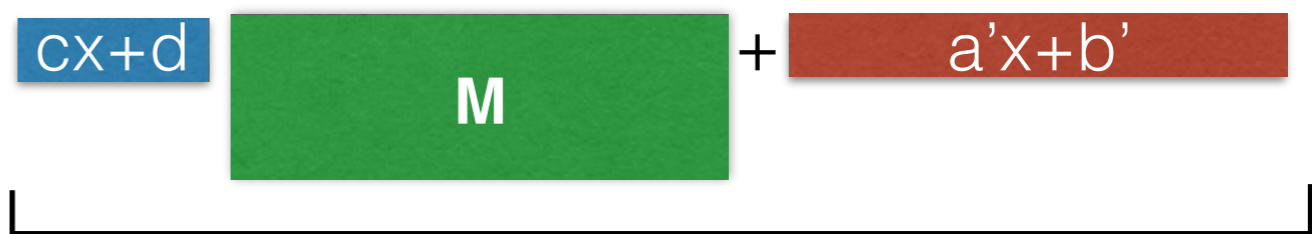
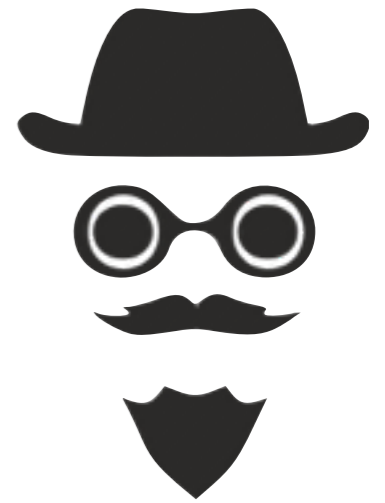
Our Method



a



b



$ax+b$

Sparse Vector-OLE Generator

2. Multi-point function secret sharing \implies sparse vector-OLE generator

- Function Secret Sharing secretly shares a function between two parties
- FSS is hard in general
- For restricted functions (point functions), it can be done very efficiently from symmetric primitives [GI14, BGI15, BGI16]
- FSS for k -point functions is basically k -sparse VOLE-generator

Improvements

- Arbitrary stretch from dual-LPN
- Linear-time evaluation from batch codes
- Improvements from LPN-friendly linear codes
- Numerous heuristic optimizations

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Efficiency (estimation): a million entries over a 128-bit field in 26 milliseconds (one core, standard laptop)

Conclusion

- Even for direct VOLE application, very competitive with alternative approaches
 - ADINZ17: $\sim 5\mu\text{s}/\text{OLE}$ (asymptotic rate 1/3)
 - Our work: $\sim 50\text{ns}/\text{OLE}$ (asymptotic rate 1/2)
- Distributed setup can be made very efficiently [Ds17]: for 1 million entries and 128-bit field, 69Kb and 34ms
- Unique communication pattern (local decompression)
- Large number of applications

Thank you for your attention

Questions?

